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**CRC CALCULATION FOR MODE-S
TRANSPONDERS**

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CRC CALCULATION FOR MODE-S TRANSPONDERS

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Summary

During Mode-S Transponder tests, it was found that the software calculation of the CRC was taking too long and degrading the performance for other essential processes. This note describes the method used for implementing a faster CRC algorithm.

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1. GENERAL

In general the Cyclic Redundancy Check is made by hardware circuitry, as shown in *Figures 2.9 and 2.10*. This is fast as it is calculated during transmission/reception of the signals and requires negligible extra time.

However, in order to verify the correct functioning of a transponder it is required to make an internal self test of the transmitted signals at regular intervals and to flag the transponder as unserviceable if this check fails.

The method used for the transponder is to verify the actual RF transmission signal against the expected signal.

The method used by the transponder in question is to generate a squitter interrogation and to count the number of low and high pulse transmissions in the outgoing pulse train and compare this with the expected number of transitions.

In order to do this, it is necessary to calculate the CRC pattern expected for the given pulse train. In this case, it is required to calculate the expected CRC by software in order to check the number of hardware transitions, all other data in the pulse train being known in the transponder.

As the transponder squitters an unsolicited DF11 interrogation at a random time interval between 800 and 1200 ms, the self test is performed on this transmission during transponder operation.

The squitter is used for the self test as it is an unsolicited transmission at regular intervals and is always active even when there is no other transponder activity.

2. PARITY CHECK CODING IN MODE-S

Parity check coding is used in Mode-S interrogations and replies to provide protection against errors. The Mode-S parity check code is defined in ICAO Annex 10 and recalled in paragraphs a), b) and c) below.

On receipt of a Mode-S interrogation, the Mode-S transponder shall perform a parity check, which is an examination of the sequence of demodulated bits to determine whether it is consistent with the code structure. If the bit sequence is consistent, the parity check is passed; otherwise it is failed. In this event, the interrogation shall not be accepted.

Similarly, the Mode-S transponder shall encode prior to transmitting a reply or squitter so that these transmissions are consistent with the code structure.

a) Parity Check Sequence :

A sequence of 24 parity check bits, generated by a code described in b) is incorporated into the field formed by the last 24 bits of all Mode-S transmissions. The 24 parity check bits are combined with either the address or the interrogation identification as described in c). The resulting combination then forms either the AP (Address / Parity) or the PI (Parity / Identification) field.

b) Parity Check Sequence Generation :

The sequence of 24 parity bits (P_1, P_2, \dots, P_{24}) is generated from the sequence of information bits (m_1, m_2, \dots, m_k) where k is 32 or 88 for short or long transmissions respectively. This is done by means of a code generated by the polynomial :

$$G(x) = \sum_{i=0}^{24} g_i x^i$$

where $g_i = 1$ for $i = 0, 3, 10$ and 12 through 24 and is 0 otherwise.

When by the application of binary polynomial algebra the above $G(x)$ is divided into $[M(x)]x^{24}$ where the information sequence $M(x)$ is expressed as :

$$M(x) = m_k + m_{k-1}x + m_{k-2}x^2 + \dots + m_1x^{k-1}$$

the result is a quotient and a remainder $R(x)$ of degree < 24 . The bit sequence formed by this remainder shall be the parity check sequence. Parity bit p_i for any i from 1 to 24, is the coefficient of x^{24-i} in $R(x)$.

Note : The effect of multiplying $M(x)$ by x^{24} is to append 24 ZERO bits to the end of the sequence.

c) AP or PI Field Generation ::

The address used for AP field generation is either the discrete address or the All-Call address. The address shall be a sequence of 24 bits, $(a_1, a_2 \dots a_{24})$. In the discrete address, a_1 shall be the bit transmitted first in the AA field of an All-Call reply. This address sequence shall be used in the downlink Address/Parity field generation, while a modified form of this sequence $(b_1, b_2 \dots b_{24})$ shall be used for uplink Address/Parity field generation.

The interrogator identifier used for PI field generation is formed by a sequence of 24 bits $(a_1, a_2 \dots a_{24})$ where the first 20 bits have zero value and the last four bits are a replica of the II field.

A modified sequence $(b_1, b_2 \dots b_{24})$ shall be used for uplink AP field generation. Bit b_i is the coefficient of x^{i-1} in the polynomial $H(x)A(x)$, where :

$$H(x) = 1 + x + x^2 + x^3 + x^4 + x^5 + x^6 + x^7 + x^8 + x^9 + x^{10} + x^{11} + x^{12} + x^{14} + x^{21} + x^{24}$$

and

$$A(x) = a_1 + a_2x + a_3x^2 + \dots + a_{24}x^{23}$$

In the SSR Mode-S address, a_i shall be the i -th bit transmitted in the AA field of an all-call. In the all-call and broadcast addresses a_i shall equal one for all values of i .

The sequence bits transmitted in the AP or PI field is :

$$t_{k+1}, t_{k+2} \dots t_{k+24}$$

The bits are numbered in order of transmission, starting with $k+1$.

In interrogations :

$$t_{k+i} = b_i \oplus P_i$$

where « \oplus » prescribes modulo-2 addition; $i=1$ is the first bit transmitted in the AP field.

In replies and squitters :

$$t_{k+i} = a_i \oplus P_i$$

where « \oplus » prescribes modulo-2 addition; $i=1$ is the first bit transmitted in the AP or PI field.

Note : *Figures Nos. 2.9 and 2.10 show typical implementation of the error protection circuits.*

3. CALCULATION OF PI FOR A SQUITTER TYPE DF 11

It can be seen from Section 2 that the PI field to be calculated for a squitter transmission consists of the application of the encoding algorithm to the data stream consisting of the DF code, the CA field, the transponder Mode-S address and the interrogator identity 0.

Hence the data stream to be encoded is :

1				32	33		56
DF	0	CA	MODE-S ADDRESS		0		0

and the transmitted data shall be :

1				32	33		56
DF	0	CA	MODE-S ADDRESS			PI	

where :

- PI is the Parity/Identification Field with $II = 0$,
- DF is the downlink Format $DF = 11$,
- CA is the capacity Field 0 to 3.

It can be seen that this particular case it is only required to pass the first 32 bits (1 to 32) through the encoding algorithm as the last 24 bits are all zero and will have no effect on the resulting PI generated.

4. METHOD USED

The method used is to divide the incoming data stream by the encoding polynomial defined for Mode-S CRC calculations.

This is done using modulo 2 addition as described in Section 2 which is equivalent to an exclusive OR function available in most computer languages.

The number of iterations required is equal to the number of bits in the incoming data stream.

Hence the execution time will depend on the number of bits in the data stream to be processed.

With long data streams methods exist where the parity may be calculated faster using look-up table(s) precalculated using the encoding polynomial as divisor but these methods are not described here.

Two examples are given of programmes to calculate a typical squitter CRC, one in PASCAL using 32 bit long word variables and the other in a MOTOROLA 6809 byte oriented machine in ASSEMBLER.

For information, the methods used for calculating the CRCs for the three normal cases used in Mode-S transmission are shown in Examples 3, 4 and 5. They are :

- a) Sensor Encoder,
- b) Transponder Decoder,
- c) Transponder Encoder or Sensor Decoder

In these examples, it is assumed that :

- The coding polynomial is FFFA0480 hex $G(x)$ of 2b),
- The message data is 20000000 hex,
- The A/C Mode-S address is DAB505 hex.

Hence the Sensor A/P field (uplink) will be 1347AC hex and the Transponder A/P field (downlink) will be 5AD35A hex.

5. CONCLUSION

The CRC calculation time in the transponder concerned was reduced from more than 40 ms to 2.4 ms, thus enabling the other processes to function normally in this real time environment.

This extremely simple method is fast enough for CRC calculations if the incoming stream is fairly short.

In any case, it is simple enough to make timing calculations depending on the number of iterations and to see if this method is suitable for a particular application.

6. REFERENCE

Minimum Operational Performance Standards for Air Traffic Control Radar Beacon System / Mode Select (ATCRBS / Mode-S) Airborne Equipment

Document No. RTCA/DO-181A

January 1992

Prepared by : SC-142

```

; CRC calculation for Mode-S transponder
; using mathematical method
    DEFSEG RAM,ABSOLUTE,CLASS=DATA
    DEFSEG ROM,ABSOLUTE,CLASS=CODE
    SEG      ROM
    ORG      $100

*
START:
    LDA      #$80      ; initialise data
    STA      DAT        ; for test only
    CLR      DAT+1
    CLR      DAT+2
    CLR      DAT+3
    CLR      J          ; clear for bit count of 32
C001:
    EQU      $
    LDD      DAT        ; Data bytes in
    BITA     #$80      ; test msb
    BEQ      C002      ; do nothing if zero
    EORA     P          ; else XOR with polynomial
    EORB     P+1
    STD      DAT
    LDD      DAT+2
    EORA     P+2
    EORB     P+3
    STD      DAT+2
C002:
    EQU      $
    LSL      DAT+3      ; shift data left 1 bit
    ROL      DAT+2
    ROL      DAT+1
    ROL      DAT

*
    LDA      J
    INC      J
    CMPA     #31      ; 32 bits ?
    BLT      C001      ; no
* result in dat
    SEG      RAM
    ORG      $200

*
P      DB      $FF,$FA,$04,$80      ; Parity polynomial
DAT    DS      4                    ; Data stream 32 bits
*
J      DS      1                    ; 32 bit count
END

```

Example 1

```

Program   CRC1 ;      { Mode-S CRC calculated by division }

{ This routine calculates the Downlink Parity Identification }
{ For a 56 bit Squitter Transmission }
{ This is done by hardware in the Sensor }
{ But verified by software in the auto-test procedure }

Uses Dos,Crt ;

VAR   j      : Integer ;

CONST
  poly : longint = ($FFFA0480) ; { Coding Polynomial }
  a : longint = ($200000000) ; { Uplink data }
  pi : longint = ($000000000) ; { Parity Identifier }

BEGIN
  { calculate Downlink Parity Identifier }
  pi := a ;
  For j := 1 to 32 do begin
    If (pi AND $800000000) <> 0 then pi := pi XOR poly ;
    pi := pi SHL 1 ;
  end ;
  { result in pi }
  { total transmission = a,pi }
END.

```

Example 2

```

Program   CRC4 ; { Mode-S ADDRESS  calculated by division }

{ This is the Transponder Encoder and Sensor Decoder }
{ This routine calculates the A/C Address or Address Parity }
{ This is done by hardware in the Transponder or Sensor }

Uses Dos,Crt ;

VAR   j           : Integer ;
        ad          : longint ;

CONST
  poly : longint  = ($FFFa0480) ; { Coding Polynomial }
  a    : longint  = ($200000000) ; { Uplink data }
  adr  : longint  = ($5ad35a00) ; { A/C Address or AP }

BEGIN

  { calculate the message parity }

  For j := 1 to 32 do begin
    If (a AND $800000000) <> 0 then a := a XOR poly ;
    { Shift total message left }
    a := a SHL 1 ;
    end ;
    { Now XOR in the A/C address or Address Parity }

    a := a XOR adr ;
    { a = aircraft address or address parity }

END.

```

Example 3

```

Program CRC3 ;      { Mode-S ADDRESS  calculated by division }

{ This is the Transponder Decoder }
{ This routine calculates the A/C Address  }
{ This is done by hardware in the transponder  }

Uses Dos,Crt ;

VAR          j          : Integer ;
              ad          : longint ;

CONST
  p : longint = ($FFFa0480) ; { Coding Polynomial }
  a : longint = ($200000000) ; { Uplink data }
  adr : longint = ($1347ac00) ; { Address Parity from Sensor }

BEGIN
{ calculate Aircraft Address  }
      ad := 0 ;
      { Treat Total Uplink Message }
      For j := 1 to 56 do begin
        If (a AND $800000000) <> 0 then a := a XOR p ;
        { Shift total message left }
        a := a SHL 1 + adr SHR 31 AND 1 ;
        adr := adr SHL 1 ;
        { for the last 24 bits }
        If j > 31 then begin
          ad := ad + a SHR 31 AND 1 ;
          ad := ad SHL 1 ;
          end ;
        end ;
      { adjust result in ad }
      ad := ad SHR 2 ;

END.

```

Example 4

```

Program CRC ;      { Mode-S CRC calculated by division }

{ This routine is the Sensor Encoder      }
{ This is done by hardware in the Sensor }

Uses Dos,Crt ;

VAR   j                : Integer ;
        par,b,ad,p       : longint ;

CONST
poly : longint = ($FFFA0480) ; { Coding Polynomial }
a : longint = ($20000000) ; { Uplink data 32 bits }
adr : longint = ($dab50500) ; { Aircraft address 24 ms bits}

BEGIN
        p := poly ;
        { 1st. calculate Uplink Message Parity }
        For j := 1 to 32 do begin
            If (a AND $80000000) <> 0 then a := a XOR poly ;
            a := a SHL 1 ;
            end ;
        { a = Uplink Message Parity }
        { Then encode the Mode-S address }
        For j := 1 to 24 do begin
            If (adr AND $80000000) <> 0 then b := b XOR p ;
            p := p SHR 1 ;
            adr := adr SHL 1 ;
            end ;
        { b = Encoded Mode-S address 24 bits }
        { Now XOR for the Uplink address parity }
        par := a XOR b ;
        { par = Total Uplink Parity 24 bits }
        { Total Uplink Message = a,par 56 bits }
END.

```

Example 5

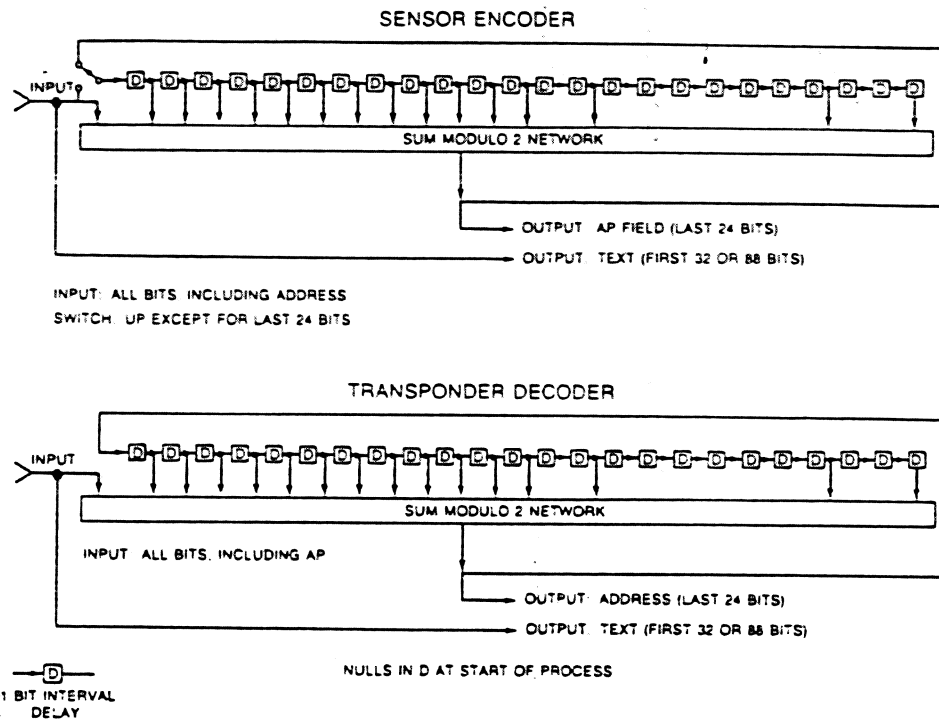


Figure 2-9 Functional Diagram of Uplink Coding

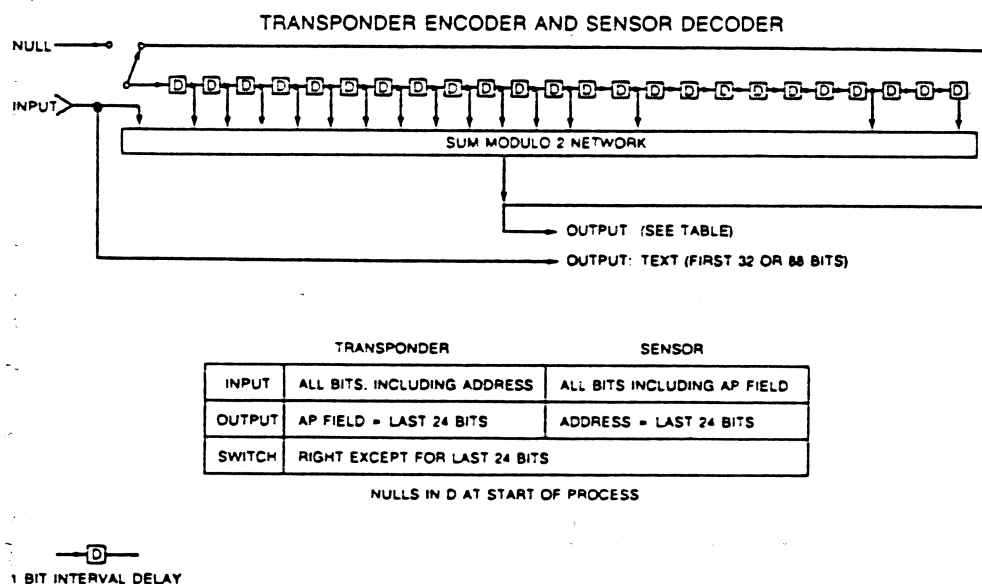


Figure 2-10 Functional Diagram of Downlink Coding